

Five Worlds: a Mathematical Problem Determines Which One We Live In

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Computation: every process which is a **sequence of simple steps**, that we want to perform, or understand

There is essential computational component in a variety of natural phenomena.

Computation in math

Solving equations: $x^2 + y^2 = z^2 \Rightarrow$ Solution $x = 3, y = 4, z = 5$

Proving theorems: $x^n + y^n = z^n, n > 2 \Rightarrow$ **NO solution, Great Fermat Theorem proved by Wiles** 200 pages paper



Can we prove theorems automatically?



Fetal development - Nature computes! In nature a computation is an evolution of an environment via repeated application of simple rules or processes.

FINDINGS

Hitting It Off, Thanks to Algorithms of Love



Computation

Mathematics

$$X^n + Y^n = Z^n$$

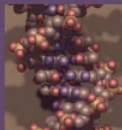
Computer Science



Computation



Physics



Biology

Plan

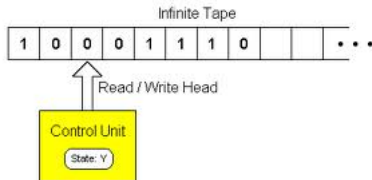
- We live in the universe where computation plays an important role.
- From Google to Genomics, the notion of an algorithm underlies this universe. **Algorithm** is the language of computation.
- Fast algorithms: **P**
- Fast verification: **NP**
- Is $P=NP$? (Is to find harder than to verify?)
- Hard problems are useful: Public key cryptography
- Five worlds of Russell Impagliazzo

Algorithm

Alan Turing 1912-1954

Formal definition of **algorithm** (Turing machine = computer with infinite memory)

Church-Turing Thesis: everything that can be computed in nature can be done on a Turing machine.



Limits on the power of algorithms

There is no algorithm to decide if a given computer program P halt on all inputs

There is no algorithm to decide if a given integer equation has an integer solution

Interesting to know

Alan Turing was the key player in the breaking of the German Enigma code at Bletchley Park during World War II.



Efficiency of an algorithm: number of basic steps for larger and larger inputs

Example: Rubik's cube



$n=2$



$n=3$



$n=4$



$n=5$

Efficiency of the addition algorithm

5 DIGITS
30 STEPS

1. Add digits. Write answer, retain carry.
2. Move one column left, write carry.
3. Scan column. If empty, stop.
4. Go to 1

12345
+6789

10 DIGITS
60 STEPS

123456789
+987654321

6 basic steps per column

20 DIGITS
120 STEPS

72635273545786043726
+53827484732625435473

50 DIGITS
300 STEPS

47563739203487456438992305757328576452364568456465744576
+98656092843467546234868431987543210979832865874134653472

N DIGITS
6N STEPS

Very fast: almost as fast as reading the input data

Efficiency of the multiplication algorithm

5 DIGITS

25 STEPS

10 DIGITS

100 STEPS

20 DIGITS

400 STEPS

50 DIGITS

2500 STEPS

N DIGITS

N^2 STEPS

Grade-school multiply algorithm

n^2

12345

x6789

123456789

x987654321

72635273545786043726

x53827484732625435473

47563739203487456438992305757328576452364568456465744576

98656092843467546234868431987543210979832865874134653472

Fast! Multiply 10,000 digits in a second!

Is there a faster algorithm?

Efficiency of a factoring algorithm

$$? \times ? = 147,573,952,588,676,412,927$$

Find nontrivial factors of a number A

N DIGITS
 10^{N^2} STEPS

Brute force factoring algorithm

Input: A

- For $B = 2, 3, \dots, \sqrt{A}$ do:
- If B divides A, return B, A/B

Very slow! 1000 digits \rightarrow sun will die before finishing

Is there a faster algorithm?

Yes, but still extremely slow!

Addition and multiplication is easy!

Is factoring hard? slow?



We don't know

The class P is the class of all problems having an efficient algorithm to find solutions

(efficient algorithm is a polynomial time algorithm, namely the number of steps to do computation for N digit numbers is bounded by N or N^2 or N^3 etc.)

Many interesting problems are in P

Addition and multiplication algorithms are in P

We will mention other examples of efficient algorithms:

Efficient algorithms: Error correction codes

Reed-Solomon decoding, 60

Petersen 60

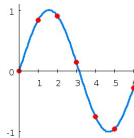
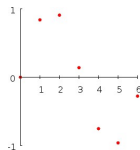
Berlecamp-Massey 68

CDs, DVDs

Satellite communication

Cell phones

Paper bar codes (printed at home tickets etc.)



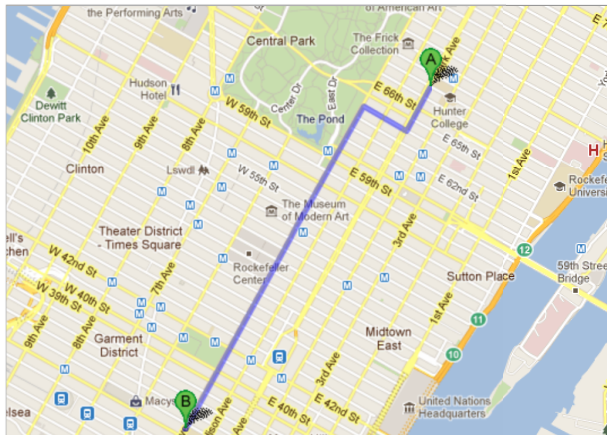
Efficient algorithms: Shortest path

Dijkstra, 1959

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To see all the details that are visible on the screen, use the "Print" link next to the map.



Efficient algorithms: Pattern matching

Rabin-Karp string search algorithm

Knuth-Morris-Pratt algorithm, 1977

Boyer-Moore string search algorithm, 1977

Text processing, Spell checking

Web search



Genome, Molecular biology

Efficient algorithms changed our life

YOU KNOW YOU ARE LIVING IN 2011 when...

1. You accidentally enter your PIN on the microwave.
2. You e-mail the person who works at the desk next to you.
3. You pull up in your own driveway and use your cell phone to see if anyone is home to help you carry in the groceries...
4. Every commercial on television has a web site at the bottom of the screen
5. Leaving the house without your cell phone, which you didn't even have 10 years ago, is now a cause for panic and you turn around to go and get it
6. You get up in the morning and go on line before getting your coffee
7. You start tilting your head sideways to smile. :)

The class NP (non-deterministic polynomial)

Are all interesting problems in P?

Cook, Levin 1971, and earlier, Godel 1956

The class NP is the class of all problems having efficient verification algorithms of given solutions

For each such problem, it is, maybe, hard to find a solution, but if somebody gives you a solution it takes polynomial time to verify that it is, indeed, a solution.

The class NP (non-deterministic polynomial)

Examples of problems in NP:

1. All problems that are in P;
2. Can a map be colored by three colors?
3. Sudoku game.

For each such problem, finding a solution (of length n) takes $\leq 2^n$ steps: try all possible solutions and verify each.

Can we do better? Do all NP problems have efficient algorithms?

Is $P=NP$?

Conjecture: $P \neq NP$.

A problem is **NP-complete** if finding a polynomial algorithm for its solution would imply $P = NP$. (Factoring is NOT NP-complete)

Are hard problems (like factoring) useful?

YES! Hard problems are useful in public key cryptography

During the early history of cryptography, two parties would rely upon a key using a secure, but non-cryptographic, method; for example, a face-to-face meeting or an exchange via a trusted courier. This key, which both parties kept absolutely secret, could then be used to exchange encrypted messages. A number of significant practical difficulties arise in this approach to distributing keys.



Radio

In 1909, Marconi and Karl Ferdinand Braun were awarded the Nobel Prize in Physics for "contributions to the development of wireless telegraphy".



⇒ **A need for Public key cryptography**

Public-key cryptography: users can communicate securely over a public channel without having to agree upon a shared key beforehand.

Diffie, Hellman, Merkle (1974)



Public key cryptography

Alice and Bob publish number g (a generator)

Alice has secret number a , Bob has secret number b .

Alice publishes g^a , Bob publishes g^b .

Now they have a secret key $g^{ab} = (g^a)^b = (g^b)^a$.

The rest of the world (Eve) has to find a or b , $a = \log_g(g^a)$ but it takes long time to compute logarithm.

Diffie-Hellman Key Exchange



Alice

Bob and Alice know and have the following :
 $p = 23$ (a prime number) $g = 11$ (a generator)



Bob

Alice chooses a secret random number $a = 6$

Alice computes : $A = g^a \bmod p$
 $A = 11^6 \bmod 23 = 9$

Alice receives $B = 5$ from Bob

Secret Key = $K = B^a \bmod p$

$$K = 5^6 \bmod 23 = 8$$

Bob chooses a secret random number $b = 5$

Bob computes : $B = g^b \bmod p$
 $B = 11^5 \bmod 23 = 5$

Bob receives $A = 9$ from Alice

Secret Key = $K = A^b \bmod p$

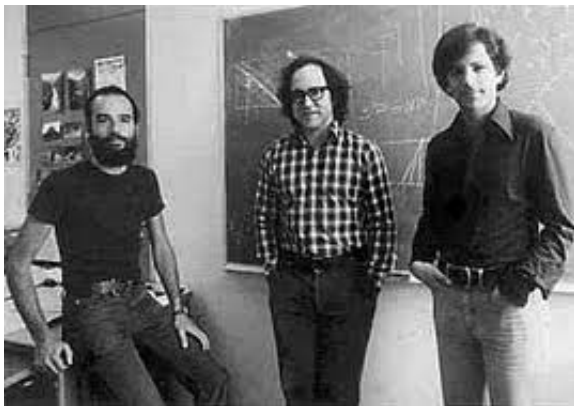
$$K = 9^5 \bmod 23 = 8$$

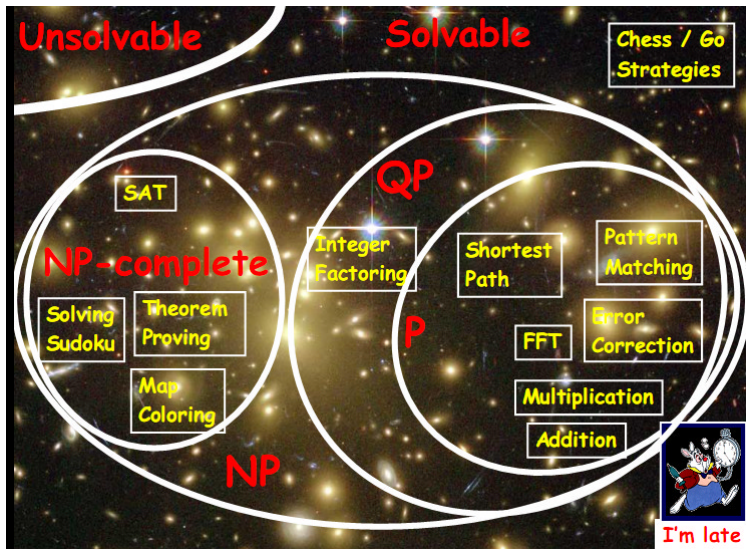
The common secret key is : 8

N.B. We could also have written : $K = g^{ab} \bmod p$

Public key cryptography

Rivest, Shamir, and Adleman (RSA)





Five worlds

Russell Impagliazzo wrote a survey on Average-Case Complexity describing 5 possible worlds: we live in one of them, but do not yet know which one.

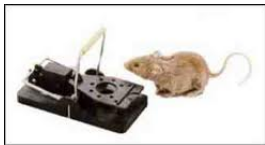
Algorithmica is the world in which $P = NP$. Utopia! Lots of science, math, engineering problems can be automated. No cryptography, no secrets, no internet exams.

If $P=NP$ we have fast, automated finder.

Heuristica is the world where NP problems are intractable in the worst case but tractable on average, the same Utopia!

Five worlds

$y = f(x)$ is **one way function** if given x it is easy to find y , but knowing y is hard to find x . Multiplication, Discrete logarithm are, probably, such functions.



Pessiland is the world in which there are hard average-case problems, but no one-way functions. This is the worst possible world.

Five worlds

Minicrypt is the world in which one-way functions exist, but public-key cryptography is impossible.

Cryptomania is the world in which public-key cryptography is possible. Electronic commerce, secure e-mail, on-line exams. Hopefully, we live in this world.



Millenium Prize Problems by Clay Math Institute

P versus NP problem

Hodge conjecture

Riemann hypothesis

Yang-Mills existence and mass gap

Navier-Stokes existence and smoothness

Birch and Swinnerton-Dyer conjecture

Poincaré conjecture (solved!)

